

Rank and nullity of Partition Regular Matrices

Leslie Hogben* Jillian McLeod†

September 6, 2009

1 Introduction

Terminology:

- Let $A \in \mathbb{Q}^{u \times v}$. We refer to A as being $CC(m)$ if A satisfies the columns condition with a partition consistent of m classes.
- By Rado's theorem we have that A is $CC(m)$ for some $m \in \mathbb{N}$ if and only if A is *partition regular*. We will use the abbreviation PR for partition regular and use it interchangeably with the statement “ A is $CC(m)$ for some $m \in \mathbb{N}$.”
- Let $A \in \mathbb{Q}^{u \times v}$. An index k is a *null index* if for every vector $\mathbf{x} = [x_i] \in \ker A$, $x_k = 0$.

Observation 1.1.

1. $A \in CC(1)$ if and only if $A\mathbf{1} = \mathbf{0}$, where $\mathbf{1} = [1, \dots, 1]^T$.
2. If A is $CC(m)$, then $\text{rank}(A) \leq v - m$, since in this case A has at least m dependent columns.
3. If $\vec{x} \in \ker(A)$ such that no coordinate of \vec{x} is allowed to be 0, and $D = \text{diag}(x_1, \dots, x_v)$, then the matrices $D^{-1}AD$ and AD are $CC(1)$ and therefore PR .

Theorem 1.2. *If A is a $u \times v$ matrix such that A has a null index, then A is not PR .*

*Department of Mathematics, Iowa State University, Ames, IA 50011, USA (lhogben@iastate.edu) and American Institute of Mathematics, 360 Portage Ave, Palo Alto, CA 94306 (hogben@aimath.org).

†Department of Mathematics, Mt. Holyoke College

2 Incidence matrices of oriented graphs

Adjacency matrices are not PR since they are nonnegative and nonzero matrices. Oriented vertex-edge incidence matrices can be PR and we look here at questions that arise naturally in this context.

Notation:

Let $\vec{G} = (V, E)$ be an oriented graph. Then $D_{\vec{G}}$ denotes its vertex-edge incidence matrix.

Observation 2.1. *While an arbitrary matrix A is $CC(1)$ if and only if $\vec{1} \in \ker(A)$, for any oriented graph \vec{G} , the vector $\mathbb{1} = [1, \dots, 1]^T$ is in the left null space of $D_{\vec{G}}$.*

Observation 2.2. *Let \vec{G} be an oriented graph and \vec{D}_G its vertex-edge incidence matrix. If \vec{G} has either a source or a sink then \vec{D}_G is not PR.*

Observation 2.3. *For any \vec{G} , the matrix $D_{\vec{G}}$ has net column weight of 0. Note that the row weight is variable.*

Theorem 2.4. *Let G be a connected graph. The following are equivalent:*

1. $K_e(G) \geq 2$, i.e., G has no bridge,
2. G is the union of its cycles,
3. G can be oriented so that the corresponding vertex-edge incidence matrix is PR.

Proof. The following recursive algorithm produces the desired partition: Pick an unoriented cycle and orient it cyclically. Let the corresponding columns of the vertex-edge incidence matrix be the first cell of our partition I_1 . Repeat this process until all edges are in some class I_k . \square

Theorem 2.5. *For an oriented graph \vec{G} , the matrix \vec{D}_G is PR if and only if \vec{G} is strongly connected.*

Corollary 2.6. *For any $D_{\vec{G}}$, where \vec{G} is strongly connected,*

$$\text{rank } D_{\vec{G}} \leq |G| - 1.$$

Observation 2.7. *If \vec{C}_n is an oriented cycle on n vertices then $D_{\vec{C}_n}$ is $CC(1)$ and $\text{rank } D_{\vec{C}_n} = n - 1$.*

Theorem 2.8. *Let G be any graph which contains a Hamiltonian cycle. Then G can be oriented in such a way that $D_{\vec{G}}$ is $CC(2)$ and therefore PR.*

Theorem 2.9. *If \vec{D}_G^1 and \vec{D}_G^2 are two distinct orientated vertex-edge incidence matrices associated with a graph G , then there exists a signature matrix S such that $\vec{D}_G^1 = \vec{D}_G^2 \cdot S$.*

Observation 2.10. *Let $D_{\vec{G}}$ be an incidence matrix for the strongly connected oriented graph \vec{G} . If $\mathbb{I} = \{I_1, \dots, I_k\}$ is a partition of the columns of $D_{\vec{G}}$ which satisfies the columns condition, then I_1 is an edge-disjoint union of cycles.*

Theorem 2.11. *Let $D_{\vec{G}}$ be an incidence matrix for the strongly connected oriented graph \vec{G} . Any set $\{I_1, \dots, I_t\}$ such that $I_i \in E_{\vec{G}}$, $I_j \cap I_l = \emptyset$ if $j \neq l$, and for all $1 \leq j \leq t$, $\sum_{i \in I_j} a_i \in \langle a_i : i \in \cup_{l=1}^t I_l \rangle$, can be extended to a partition of $E_{\vec{G}}$ that satisfies the columns condition.*

Corollary 2.12. *Let $D_{\vec{G}}$ be an incidence matrix for the strongly connected oriented graph \vec{G} . The greedy algorithm produces a partition of the columns of $D_{\vec{G}}$ which satisfies the columns condition.*

3 Sign Patterns

A *sign pattern matrix* (or *sign pattern* for short) is a (rectangular) matrix having entries in $\{+, -, 0\}$. For a real matrix A , $\text{sgn}(A)$ is the sign pattern having entries that are the signs of the corresponding entries in A . If \mathbb{Y} is an $n \times n$ sign pattern, the *sign pattern class* (or *qualitative class*) of \mathbb{Y} , denoted $\mathcal{Q}(\mathbb{Y})$, is the set of all $A \in \mathbb{R}^{n \times n}$ such that $\text{sgn}(A) = \mathbb{Y}$. It is traditional in the study of sign patterns to say that a sign pattern \mathbb{Y} *requires* property P if every matrix in $\mathcal{Q}(\mathbb{Y})$ has property P and to say that \mathbb{Y} *allows* property P if there exists a matrix in $\mathcal{Q}(\mathbb{Y})$ that has property P . Patterns that require partition regularity are too trivial to be of interest:

Theorem 3.1. *The only sign patterns that requires partition regularity are the all zero sign patterns.*

Proof. Assume $\mathbb{Y} = [\psi_{ij}]$ has a nonzero entry. Construct a matrix $A = [a_{ij}]$ as follows:

- For all i, j such that $\psi_{ij} = 0$, $a_{ij} = 0$.
- For all i, j such that $\psi_{ij} = +$, $a_{ij} = 1$.
- For all i, j such that $\psi_{ij} = -$, $a_{ij} = -\frac{1}{n}$.

There is no subset of columns that sum to zero, so A does not have the columns condition and so is not partition regular. \square

Since a partition regular matrix must satisfy the columns condition, it is clear that in order for a sign pattern to allow partition regularity, any nonzero row must have both at least one $+$ entry and at least one $-$ entry. This is also sufficient for a sign pattern to allow partition regularity:

Theorem 3.2. *Let \mathbb{Y} be an $m \times n$ sign pattern. The following are equivalent:*

1. *Each row of \mathbb{Y} has at least one $+$ entry and at least one $-$ entry or every entry is 0.*
2. *\mathbb{Y} allows CC(1).*
3. *\mathbb{Y} allows partition regularity.*

Proof. It is clear that (2) \implies (3) \implies (1). Assume each row of $\mathbb{Y} = [\psi_{ij}]$ has at least one $+$ entry and at least one $-$ entry or every entry is 0. If row i is non entirely 0, let $m(i)$ denote the column number of the first $-$ entry in row i ; otherwise, $m(i) = 0$. Construct a matrix $A = [a_{ij}]$ as follows:

- For all i, j such that $\psi_{ij} = 0$, $a_{ij} = 0$.
- For all i such that $m(i) > 0$:
 - If $\psi_{ij} = +$, $a_{ij} = 1$.
 - For $j > m(i)$, if $\psi_{ij} = -$ then $a_{ij} = -\frac{1}{n}$.
 - $a_{i, m(i)} = -\sum_{j \neq m(i)} a_{ij}$.

Clearly $A \in \mathcal{Q}(\mathbb{Y})$ and $A\mathbf{1} = \mathbf{0}$, so A has CC(1). \square

The *minimum rank* of an $m \times n$ sign pattern \mathbb{Y} is

$$\text{mr}(\mathbb{Y}) = \min\{\text{rank}(A) : A \in \mathcal{Q}(\mathbb{Y})\},$$

and the *maximum nullity* of \mathbb{Y} is

$$\text{M}(\mathbb{Y}) = \max\{\text{null}(A) : A \in \mathcal{Q}(\mathbb{Y})\}.$$

Clearly $\text{mr}(\mathcal{Y}) + \text{M}(\mathcal{Y}) = n$.

It is not always the case that the nullity of a partition regular matrix can be realized by the number of cells in a columns condition partition. For example, for $A = \begin{bmatrix} 3 & -1 & -1 & -1 \end{bmatrix}$, $\text{null}(A) = 3$ but A is $\text{CC}(m)$ only for $m = 1$. But the following question remains open:

Question 3.3. *If \mathbb{Y} allows partition regularity, must there exist a matrix $A \in \mathcal{Q}(\mathbb{Y})$ such that A is $\text{CC}(\text{M}(\mathbb{Y}))$?*

Theorem 3.4. *Let G be a connected graph, let \vec{G} be an orientation of G , and let $\mathbb{Y} = \text{sgn}(D_{\vec{G}})$. Then the following are equivalent:*

1. \mathbb{Y} allows PR
2. $D_{\vec{G}}$ is PR
3. \vec{G} is strongly connected.

Proof. [will write this later]

□